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On the Linear Convergence Rate of the Distributed Block Proximal Method

Francesco Farina, *Member, IEEE*, Giuseppe Notarstefano, *Member, IEEE*

Abstract—The recently developed Distributed Block Proximal Method, for solving stochastic big-data convex optimization problems, is studied in this paper under the assumption of constant stepsizes and strongly convex (possibly non-smooth) local objective functions. This class of problems arises in many learning and classification problems in which, for example, strongly-convex regularizing functions are included in the objective function, the decision variable is extremely high dimensional, and large datasets are employed. The algorithm produces local estimates by means of block-wise updates and communication among the agents. The expected distance from the (global) optimum, in terms of cost value, is shown to decay linearly to a constant value which is proportional to the selected local stepsizes. A numerical example involving a classification problem corroborates the theoretical results.

Index Terms—Optimization algorithms, Machine learning, Network analysis and control

I. INTRODUCTION

In this paper, we address in a distributed way stochastic big-data convex optimization problems involving *strongly convex* (possibly nonsmooth) local objective functions, by means of the Distributed Block Proximal Method [1], [2]. Problems with this structure naturally arise in many learning and control problems in which the decision variable is extremely high dimensional and large datasets are employed. Relevant examples include: direct policy search in reinforcement learning [3], dynamic problems involving stochastic functions generated from collected samples to be processed online [4], learning problems involving massive datasets in which sample average approximation techniques are used [5], and settings in which only noisy subgradients of the objective functions can be computed at each time instant [6].

Distributed algorithms for solving stochastic problems have been widely studied [6]–[11]. On the other side, distributed algorithms for big-data problems through block communication have started to appear only recently [12]–[16]. The Distributed Block Proximal Method solves problems that can be together non-smooth, stochastic and big-data, thus distinguishing from the above works (see [1] for a comprehensive literature review). This algorithm evolves through block-wise communication and updates (involving subgradients of the local functions and proximal mappings induced by some distance generating functions) and it has been already shown to achieve a

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sublinear convergence rate on problems with non-smooth convex objective functions. The contribution of this paper is to extend this result by showing that, under strongly convex (possibly non-smooth) local objective functions and constant stepsizes, the Distributed Block Proximal Method exhibits a linear convergence rate (with a constant error term) to the optimal cost in expected value. The main challenge in the linear-rate analysis relies in the block-wise nature of the algorithm.

II. SET-UP AND PRELIMINARIES

A. Notation, definitions and preliminary results

Given a vector $x \in \mathbb{R}^n$, we denote by x_ℓ the ℓ -th block of x , i.e., given a partition of the identity matrix $I = [U_1, \dots, U_B]$, with $U_\ell \in \mathbb{R}^{n \times n_\ell}$ for all ℓ and $\sum_{\ell=1}^B n_\ell = n$, it holds $x = \sum_{\ell=1}^B U_\ell x_\ell$ and $x_\ell = (U_\ell)^\top x$. In general, given a vector $x_i \in \mathbb{R}^n$, we denote by $x_{i,\ell}$ the ℓ -th block of x_i . Given a matrix A , we denote by a_{ij} the element of A located at row i and column j . Given two vectors $a, b \in \mathbb{R}^n$ we denote by $\langle a, b \rangle$ their scalar product. Given a discrete random variable $r \in \{1, \dots, R\}$, we denote by $P(r = \bar{r})$ the probability of r to be equal to \bar{r} for all $\bar{r} \in \{1, \dots, R\}$. Given a nonsmooth function f , we denote by $\partial f(x)$ its subdifferential at x .

The following preliminary result will be used in the paper.

Lemma 1. *Given any two scalars $\delta \neq \gamma \neq 1$, it holds that*

- (i) $\sum_{s=r}^t \delta^s = \frac{\delta^r - \delta^{t+1}}{1 - \delta}$
- (ii) $\sum_{s=0}^t \delta^{t-s} \gamma^s = \frac{\delta^{t+1} - \gamma^{t+1}}{\delta - \gamma}$. □

B. Distributed stochastic optimization set-up

Let us start by formalizing the optimization problem addressed in this paper. We consider problems in the form

$$\underset{x \in X}{\text{minimize}} \quad \sum_{i=1}^N \mathbb{E}[h_i(x; \xi_i)]. \quad (1)$$

where ξ_i is a random variable and $x \in \mathbb{R}^n$, $n \gg 1$, has a block structure, i.e., $x = [x_1^\top, \dots, x_B^\top]^\top$, with $x_\ell \in \mathbb{R}^{n_\ell}$ for all ℓ and $\sum_{\ell=1}^B n_\ell = n$. The decision variable x can be very high-dimensional, which calls for block-wise algorithms.

Let $f_i(x) \triangleq \mathbb{E}[h_i(x; \xi_i)]$. Moreover, let $g_i(x; \xi_i) \in \partial h_i(x; \xi_i)$ (resp. $g_i(x) \in \partial f_i(x)$) be a subgradient of $h_i(x; \xi_i)$ (resp. $f_i(x)$) computed at x . Then, the following assumption holds for problem (1).

Assumption 1 (Problem structure).

- (A) The constraint set X has the block structure $X = X_1 \times \dots \times X_B$, where, for $\ell = 1, \dots, B$, the set $X_\ell \subseteq \mathbb{R}^{n_\ell}$ is closed and convex, and $\sum_{\ell=1}^B n_\ell = n$.

- (B) The function $h_i(x, \xi_i) : \mathbb{R}^n \rightarrow \mathbb{R}$ is continuous, strongly convex and possibly nonsmooth for all $x \in X$ and every ξ_i , for all $i \in \{1, \dots, N\}$. In particular, there exists a constant $m > 0$ such that $f_i(a) \geq f_i(b) - \langle g_i(b), b-a \rangle + \frac{m}{2} \|a-b\|^2$, for all $a, b \in X$ and all $i \in \{1, \dots, N\}$.
- (C) every subgradient $g_i(x; \xi_i)$ is an unbiased estimator of the subgradient of f_i , i.e., $\mathbb{E}[g_i(x; \xi_i)] = g_i(x)$. Moreover, there exist constants $G_i \in [0, \infty)$ and $\bar{G}_i \in [0, \infty)$ such that $\mathbb{E}[\|g_i(x; \xi_i)\|] \leq G_i$, and $\mathbb{E}[\|g_i(x; \xi_i)\|^2] \leq \bar{G}_i$, for all x and ξ_i , for all $i \in \{1, \dots, N\}$. \square

Let us denote by $g_{i,\ell}(x; \xi_i)$ the ℓ -th block of $g_i(x; \xi_i)$ and let $g(x) \in \partial f(x)$ be a subgradient of f computed at x . Then, Assumption 1(C) implies that $\mathbb{E}[\|g_{i,\ell}(x; \xi_i)\|] \leq G_i$ for all ℓ and $\|g_i(x)\| \leq G_i$. Moreover, let $\bar{G} \triangleq \sum_{i=1}^N \bar{G}_i$ and $G \triangleq \sum_{i=1}^N G_i$. Then, $\|g(x)\| \leq G$ and $\|g_i(x)\| \leq G$ for all i .

Problem (1) is to be cooperatively solved by a network of N agents. Agents locally know only a portion of the entire optimization problem. Namely, agent i knows only $g_i(x; \xi_i)$ for any x and ξ_i , and the constraint set X . The network communication is assumed to satisfy the next assumption.

Assumption 2 (Communication structure).

- (A) The network is modeled through a weighted *strongly connected* directed graph $\mathcal{G} = (\mathcal{V}, \mathcal{E}, W)$ with $\mathcal{V} = \{1, \dots, N\}$, $\mathcal{E} \subseteq \mathcal{V} \times \mathcal{V}$ and $W \in \mathbb{R}^{N \times N}$ being the weighted adjacency matrix. We define $\mathcal{N}_{i,out} \triangleq \{j \mid (i, j) \in \mathcal{E}\} \cup \{i\}$ and $\mathcal{N}_{i,in} \triangleq \{j \mid (j, i) \in \mathcal{E}\} \cup \{i\}$.
- (B) For all $i, j \in \{1, \dots, N\}$, the weights w_{ij} of the weight matrix W satisfy
- (i) $w_{ij} > 0$ if and only if $j \in \mathcal{N}_{i,in}$;
 - (ii) there exists a constant $\eta > 0$ such that $w_{ii} \geq \eta$ and if $w_{ij} > 0$, then $w_{ij} \geq \eta$;
 - (iii) $\sum_{j=1}^N w_{ij} = 1$ and $\sum_{i=1}^N w_{ij} = 1$. \square

In order to solve problem (1) agents will be using ad-hoc *proximal mappings* (see, e.g., [17]). In particular, a function ω_ℓ is associated to the ℓ -th block of the optimization variable for all ℓ . Let the function $\omega_\ell : X_\ell \rightarrow \mathbb{R}$, be continuously differentiable and σ_ℓ -strongly convex. Functions ω_ℓ are sometimes referred to as distance generating functions. Then, we define the *Bregman's divergence* associated to ω_ℓ as

$$\nu_\ell(a, b) = \omega_\ell(b) - \omega_\ell(a) - \langle \nabla \omega_\ell(a), b - a \rangle,$$

for all $a, b \in X_\ell$. Moreover, given $a \in X_\ell$, $b \in \mathbb{R}^{n_\ell}$ and $c \in \mathbb{R}$, the proximal mapping associated to ν_ℓ is defined as

$$\text{prox}_\ell(a, b, c) = \arg \min_{u \in X_\ell} \left(\langle b, u \rangle + \frac{1}{c} \nu_\ell(a, u) \right). \quad (2)$$

We make the following assumption on the functions ν_ℓ .

Assumption 3 (Bregman's divergences properties).

- (A) There exists a constant $Q > 0$ such that

$$\nu_\ell(a, b) \leq \frac{Q}{2} \|a - b\|^2, \quad \forall a, b \in X_\ell \quad (3)$$

for all $\ell \in \{1, \dots, B\}$.

- (B) For all $\ell \in \{1, \dots, B\}$, the function ν_ℓ satisfies

$$\nu_\ell \left(\sum_{j=1}^N \theta_j a_j, b \right) \leq \sum_{j=1}^N \theta_j \nu_\ell(a_j, b), \quad \forall a_1, \dots, a_N, b \in X_\ell, \quad (4)$$

where $\sum_{j=1}^N \theta_j = 1$ and $\theta_j \geq 0$ for all j . \square

Notice that Assumption 3(A) implies that, given any two points $a, b \in X$,

$$\sum_{\ell=1}^B \nu_\ell(a_\ell, b_\ell) \leq \frac{Q}{2} \sum_{\ell=1}^B \|a_\ell - b_\ell\|^2 = \frac{Q}{2} \|a - b\|^2. \quad (5)$$

Moreover, Assumption 3(B) is satisfied by many functions (such as the quadratic function and the exponential function) and conditions on ω_ℓ guaranteeing (4) can be provided [18].

III. DISTRIBUTED BLOCK PROXIMAL METHOD

Let us now recall the Distributed Block Proximal Method for solving problem (1) in a distributed way. The pseudocode of the algorithm is reported in Algorithm 1, where, for notational convenience, we defined $g_{i,\ell}(t) \triangleq g_{i,\ell}(y_i(t); \xi_i(t))$. We refer to [1] for all the details.

Algorithm 1 Distributed Block Proximal Method

Initialization: $x_i(0)$

Evolution: for $t = 0, 1, \dots$

UPDATE for all $j \in \mathcal{N}_{i,in}$

$$x_{j,\ell_i}(t) = \begin{cases} x_{j,\ell}(t), & \text{if } \ell = \ell_j(t-1) \text{ and } s_j(t-1) = 1 \\ x_{j,\ell_i}(t-1), & \text{otherwise} \end{cases}$$

if $s_i^t = 1$ then

PICK $\ell_i(t) \in \{1, \dots, B\}$ with $P(\ell_i(t) = \ell) = p_{i,\ell} > 0$

COMPUTE

$$y_i(t) = \sum_{j \in \mathcal{N}_{i,in}} w_{ij} x_{j|\ell_i}(t)$$

UPDATE

$$x_{i,\ell}(t+1) = \begin{cases} \text{prox}_\ell \left(y_i(t), g_{i,\ell}(t), \alpha_i \right), & \text{if } \ell = \ell_i(t) \\ x_{i,\ell}(t), & \text{otherwise} \end{cases}$$

BROADCAST $x_{i,\ell_i(t)}(t+1)$ to $j \in \mathcal{N}_{i,out}$

else $x_i(t+1) = x_i(t)$

The algorithm works as follows. Each agent i maintains a local solution estimate $x_i(t)$ and a local copy of the estimates of its in-neighbors (namely, $x_{j|\ell_i}(t)$ denotes the copy of the solution estimate of agent j at agent i). The initial conditions are initialized with random (bounded) values $x_i(0)$ which are shared between neighbors. At each iteration, agents can be awake or idle, thus modeling a possible asynchrony in the network. The probability of agent i to be awake is denoted by $p_{i,on} \in (0, 1]$. If agent i is awake at iteration t , it picks randomly a block $\ell_i(t) \in \{1, \dots, B\}$, some $\xi_i(t)$, and performs two updates:

- (i) it computes a weighted average of its in-neighbors' estimates $x_{j|\ell_i}(t)$, $j \in \mathcal{N}_{i,in}$;
- (ii) it computes $x_i(t+1)$ by updating the $\ell_i(t)$ -th block of $x_i(t)$ through a proximal mapping step (with a constant stepsize α_i) and leaving the other blocks unchanged.

Finally, it broadcasts $x_{i,\ell_i(t)}(t+1)$ to its out-neighbors. The status (awake or idle) of node i at iteration t is modeled as a random variable $s_i(t) \in \{0,1\}$ which is 1 with probability $p_{i,on}$ and 0 with probability $1 - p_{i,on}$.

As already stated in [1], it is worth remarking that all the quantities involved in the Distributed Block Proximal Method are local for each node. In fact, each node has locally defined probabilities (both of awakening and block drawing) and local stepsizes. Moreover, it is worth recalling that, from [1, Lemma 5], we have $x_{j|z}(t) = x_j(t)$ for all t and hence, Algorithm 1 can be compactly rewritten as follows. For all $i \in \{1, \dots, N\}$ and all t , if $s_i(t) = 1$,

$$y_i(t) = \sum_{j=1}^N w_{ij} x_j(t), \quad (6)$$

$$x_{i,\ell}(t+1) = \begin{cases} \text{prox}_{\ell} \left(y_{i,\ell}(t), g_{i,\ell}(t), \alpha_i \right), & \text{if } \ell = \ell_i(t), \\ x_{i,\ell}(t), & \text{otherwise,} \end{cases} \quad (7)$$

else, $x_i(t+1) = x_i(t)$. We will use (6)-(7) in place of Algorithm 1, in the following analysis.

IV. ALGORITHM ANALYSIS AND CONVERGENCE RATE

Let $\mathbf{x}(\tau) \triangleq [x_1(\tau)^\top, \dots, x_N(\tau)^\top]^\top$ and let $\mathcal{S}(t) \triangleq \{\mathbf{x}(\tau) \mid \tau \in \{0, \dots, t\}\}$ be the set of estimates generated by the Distributed Block Proximal Method up to iteration t . Moreover, define the probability of node i to both be awake and pick block ℓ as

$$\pi_{i,\ell} \triangleq p_{i,on} p_{i,\ell}$$

and define $a \triangleq [\alpha_1, \dots, \alpha_N]^\top$, $a_M \triangleq \max_i \alpha_i$ and $a_m \triangleq \min_i \alpha_i$. Moreover, define the average (over the agents) of the local estimates at t as

$$\bar{x}(t) \triangleq \frac{1}{N} \sum_{i=1}^N x_i(t). \quad (8)$$

Finally, let us make the following assumption about the random variables involved in the algorithm.

Assumption 4 (Random variables).

- (A) The random variables $\ell_i(t)$ and $s_i(t)$ are independent and identically distributed for all t , for all $i \in \{1, \dots, N\}$.
- (B) For any given t , the random variables $s_i(t)$, $\ell_i(t)$ and $\xi_i(t)$ are independent of each other for all $i \in \{1, \dots, N\}$.
- (C) There exists a constants $C_i \in [0, \infty)$ such that $\mathbb{E}[\|x_i(0)\|] \leq C_i$ for all $i \in \{1, \dots, N\}$ and hence $\mathbb{E}[\|\mathbf{x}(0)\|] \leq C = \sum_{i=1}^N C_i$. \square

In the following we analyze the convergence properties of the Distributed Block Proximal Method with constant stepsizes under the previous assumptions. We start by showing that consensus is achieved in the network, by specializing the results in [1]. Then, we show that also optimality is achieved in expected value and with a constant error, by studying the properties of an ad-hoc Lyapunov-like function. Finally, we show how the main result implies a linear convergence rate for the algorithm.

A. Reaching consensus

The following lemma characterizes the expected distance of $x_i(t)$ and $y_i(t)$ from the average $\bar{x}(t)$ (defined in (8)).

Lemma 2. *Let Assumptions 1, 2, 4 hold. Then, there exist constants $M \in (0, \infty)$ and $\mu_M \in (0, 1)$ such that*

$$\mathbb{E}[\|x_i(t) - \bar{x}(t)\|] \leq \mu_M^{t-1} \bar{R} + \bar{S}, \quad (9)$$

$$\mathbb{E}[\|y_i(t) - x_i(t)\|] \leq 2\mu_M^{t-1} \bar{R} + 2\bar{S} \quad (10)$$

for all $i \in \{1, \dots, N\}$ and all $t \geq 1$, with $\bar{R} = MB \left(C - \frac{a_M G}{\sigma(1-\mu_M)} \right)$ and $\bar{S} = a_M \frac{MBG}{\sigma} \frac{2-\mu_M}{1-\mu_M}$.

Proof. The proof follows by using constant stepsizes in [1, Lemma 7 and Lemma 8]. \square

In the next section, in order to prove the convergence to the optimal cost with a linear rate, we will need the following result assuring the boundedness of a particular quantity. In particular, given a scalar $c \in (0, 1)$, let us define

$$\beta(t) \triangleq \sum_{\tau=0}^t c^{t-\tau} \mathbb{E}[\|x_i(\tau) - \bar{x}(\tau)\|]. \quad (11)$$

Then, the next lemma provides a bound on $\beta(t)$ for all t .

Lemma 3. *Let Assumptions 1, 2, 4 hold. Then, for any scalar $c \in (0, 1)$,*

- (i) if $c \neq \mu_M$,

$$\beta(t) \leq c^t \left(C + \frac{\bar{R}}{c - \mu_M} \right) + \frac{1 - c^t}{1 - c} \bar{S} \quad (12)$$

- (ii) if $c = \mu_M$,

$$\beta(t) \leq c^t \left(C + \frac{t\bar{R}}{c} \right) + \frac{1 - c^t}{1 - c} \bar{S} \quad (13)$$

for all $i \in \{1, \dots, N\}$, for all t .

Proof. By using Assumption 4(C), for $\tau = 0$, one has

$$\begin{aligned} \mathbb{E}[\|x_i(0) - \bar{x}(0)\|] &\leq \mathbb{E}[\|x_i(0)\|] + \mathbb{E}[\|\bar{x}(0)\|] \\ &\leq C_i + \frac{1}{N} \sum_{j=1}^N C_j \leq C_i + \max_j C_j \leq C \end{aligned} \quad (14)$$

Hence, $\beta(t) \leq c^t C + \sum_{\tau=1}^t c^{t-\tau} \mathbb{E}[\|x_i(\tau) - \bar{x}(\tau)\|]$ and, from Lemma 2, we have

$$\beta(t) \leq c^t C + \bar{R} \sum_{\tau=1}^t c^{t-\tau} \mu_M^{\tau-1} + \bar{S} \sum_{\tau=1}^t c^{t-\tau} \quad (15)$$

Let us consider the case $c \neq \mu_M$. By using Lemma 1, one easily gets

$$\begin{aligned} \beta(t) &\leq c^t C + \frac{c^t - \mu_M^t}{c - \mu_M} \bar{R} + \frac{1 - c^t}{1 - c} \bar{S} \\ &\leq c^t \left(C + \frac{\bar{R}}{c - \mu_M} \right) + \frac{1 - c^t}{1 - c} \bar{S} \end{aligned}$$

where in the second line we have removed the negative term depending on μ_M^t . For the case $c = \mu_M$ we have

$$\sum_{\tau=1}^t c^{t-\tau} \mu_M^{\tau-1} = \sum_{\tau=1}^t c^{t-1} = t c^{t-1} \quad (16)$$

and (13) is obtained by substituting (16) in (15) and using Lemma 1. \square

B. Reaching optimality

Let us start by defining a Lyapunov-like function

$$V_i^\tau \triangleq \sum_{\ell=1}^B \pi_{i,\ell}^{-1} \nu_\ell(x_{i,\ell}^\tau, x_\ell^*) \quad (17)$$

and let $V^t \triangleq \sum_{i=1}^N V_i^t$. Moreover, define

$$f_{\text{best}}(\bar{x}^t) \triangleq \min_{\tau \leq t} \mathbb{E}[f(\bar{x}^\tau)] \quad (18)$$

and $\pi_m = \min_{i,\ell} \pi_{i,\ell}$. Then, the following result holds true and will be the key for proving the linear convergence rate of the Distributed Block Proximal Method under the previous assumptions.

Lemma 4. *Let Assumptions 1, 2, 3 and 4 hold. Moreover, let $\alpha_i \leq \frac{Q}{m}$ for all i . Then, for all t ,*

$$\begin{aligned} \mathbb{E}[V(t+1)] &\leq \left(1 - \frac{ma_m \pi_m}{Q}\right) \mathbb{E}[V(t)] \\ &\quad - \sum_{i=1}^N \alpha_i (\mathbb{E}[f_i(y_i(t))] - f_i(x^*)) + \frac{a_M^2 \bar{G}}{2\sigma}. \end{aligned} \quad (19)$$

Proof. In order to simplify the notation, let us denote $\mathbf{g}_i(t) = \mathbf{g}_i(y_i(t))$. By using the same arguments used in the proof of [1, Theorem 1] we have

$$\begin{aligned} \mathbb{E}[V_i(t+1) | \mathcal{S}(t)] &\leq V_i(t) - \sum_{\ell=1}^B \nu_\ell(x_{i,\ell}(t), x_\ell^*) \\ &\quad + \sum_{\ell=1}^B \nu_\ell(y_{i,\ell}(t), x_\ell^*) - \alpha_i \langle \mathbf{g}_i(t), y_i(t) - x^* \rangle + \frac{\alpha_i^2 \bar{G}_i}{2\sigma} \end{aligned} \quad (20)$$

Now, By exploiting Assumptions 1(B), 3(A) and (3), one has that, for all t ,

$$\begin{aligned} &\alpha_i \langle \mathbf{g}_i(t), y_i(t) - x^* \rangle \\ &\geq \alpha_i \left(f_i(y_i(t)) - f_i(x^*) + \frac{m}{2} \|y_i(t) - x^*\|^2 \right) \\ &\geq \alpha_i (f_i(y_i(t)) - f_i(x^*)) + \frac{m\alpha_i}{Q} \sum_{\ell=1}^B \nu_\ell(y_{i,\ell}(t), x_\ell^*). \end{aligned} \quad (21)$$

Now, by using (21) in (20) and by exploiting the fact that $\alpha_i \leq \frac{Q}{m}$, we get

$$\begin{aligned} \mathbb{E}[V_i(t+1) | \mathcal{S}(t)] &\leq V_i(t) - \sum_{\ell=1}^B \nu_\ell(x_{i,\ell}(t), x_\ell^*) \\ &\quad + \left(1 - \frac{m\alpha_i}{Q}\right) \sum_{\ell=1}^B \nu_\ell(y_{i,\ell}(t), x_\ell^*) \\ &\quad - \alpha_i (f_i(y_i(t)) - f_i(x^*)) + \frac{\alpha_i^2 \bar{G}_i}{2\sigma} \\ &\leq V_i(t) - \sum_{\ell=1}^B \nu_\ell(x_{i,\ell}(t), x_\ell^*) \\ &\quad + \left(1 - \frac{m\alpha_i}{Q}\right) \sum_{j=1}^N w_{ij} \sum_{\ell=1}^B \nu_\ell(x_{j,\ell}(t), x_\ell^*) \\ &\quad - \alpha_i (f_i(y_i(t)) - f_i(x^*)) + \frac{\alpha_i^2 \bar{G}_i}{2\sigma}, \end{aligned} \quad (22)$$

where in the second inequality we used assumption 3(B). If we now sum over i , by noticing that $a_m \leq \alpha_i$ for all i , we obtain

$$\begin{aligned} \sum_{i=1}^N \mathbb{E}[V_i^{t+1} | \mathcal{S}(t)] &\leq \sum_{i=1}^N V_i(t) - \sum_{i=1}^N \sum_{\ell=1}^B \nu_\ell(x_{i,\ell}(t), x_\ell^*) \\ &\quad + \sum_{i=1}^N \left(1 - \frac{m\alpha_i}{Q}\right) \sum_{j=1}^N w_{ij} \sum_{\ell=1}^B \nu_\ell(x_{j,\ell}(t), x_\ell^*) \\ &\quad - \sum_{i=1}^N \alpha_i (f_i(y_i(t)) - f_i(x^*)) + \sum_{i=1}^N \frac{\alpha_i^2 \bar{G}_i}{2\sigma}. \end{aligned} \quad (23)$$

Now, by using the fact that $a_m \leq \alpha_i \leq a_M$ for all i , the double stochasticity of W from Assumption 2, and the definition of \bar{G} , one easily obtains that

$$\begin{aligned} &\sum_{i=1}^N \mathbb{E}[V_i^{t+1} | \mathcal{S}(t)] \\ &\leq \sum_{i=1}^N V_i(t) - \frac{ma_m}{Q} \sum_{i=1}^N \sum_{\ell=1}^B \nu_\ell(x_{i,\ell}(t), x_\ell^*) \\ &\quad - \sum_{i=1}^N \alpha_i (f_i(y_i(t)) - f_i(x^*)) + \frac{a_M^2 \bar{G}}{2\sigma}. \end{aligned} \quad (24)$$

Moreover, by using (17) we can rewrite

$$\begin{aligned} &\sum_{i=1}^N V_i(t) - \frac{ma_m}{Q} \sum_{i=1}^N \sum_{\ell=1}^B \nu_\ell(x_{i,\ell}(t), x_\ell^*) \\ &= \sum_{i=1}^N \sum_{\ell=1}^B \left(\pi_{i,\ell}^{-1} \nu_\ell(x_{i,\ell}(t), x_\ell^*) - \frac{ma_m}{Q} \nu_\ell(x_{i,\ell}(t), x_\ell^*) \right) \\ &\leq \left(1 - \frac{m\pi_m a_m}{Q}\right) V(t) \end{aligned} \quad (25)$$

where we have used the fact that $\sum_{\ell=1}^B \pi_{i,\ell}^{-1} \nu_\ell(a, b) \geq \sum_{\ell=1}^B \pi_{i,\ell}^{-1} \nu_\ell(a, b)$. Finally, by plugging (25) in (24) and by using tower property of conditional expectation one gets (19). \square

Thanks to the previous results, we are now ready to state and prove the main result of this paper.

Theorem 1. *Let Assumptions 1, 2, 3 and 4 hold. Moreover, let $\alpha_i \leq \frac{Q}{m}$ for all i and let $c \triangleq \left(1 - \frac{ma_m \pi_m}{Q}\right)$. Then,*

1) if $c \neq \mu_M$,

$$f_{\text{best}}(\bar{x}^t) - f(x^*) \leq \frac{c^t}{1 - c^{t+1}} (Q + R_1) + S, \quad (26)$$

2) if $c = \mu_M$,

$$f_{\text{best}}(\bar{x}^t) - f(x^*) \leq \frac{c^t}{1 - c^{t+1}} (Q + tR_2) + S, \quad (27)$$

where $Q = (1 - c) \left(\frac{\mathbb{E}[V^0]}{a_m} + 3GC \right)$, $R_1 = \frac{(1-c)3G\bar{R}}{c - \mu_M}$, $R_2 = \frac{(1-c)3G\bar{R}}{c}$. and $S = \frac{a_M^2 \bar{G}}{2\sigma a_m} + 3G\bar{S}$.

Proof. By recursively applying (19), one has

$$\begin{aligned} & \sum_{\tau=0}^t c^{t-\tau} \sum_{i=1}^N \alpha_i (\mathbb{E}[f_i(y_i(\tau))] - f_i(x^*)) \\ & \leq c^{t+1} \mathbb{E}[V^0] + \sum_{\tau=0}^t c^{t-\tau} \frac{a_M^2 \bar{G}}{2\sigma} \end{aligned}$$

Moreover, since $a_m \leq \alpha_i$ for all i ,

$$\begin{aligned} & \sum_{\tau=0}^t c^{t-\tau} a_m \sum_{i=1}^N (\mathbb{E}[f_i(y_i(\tau))] - f_i(x^*)) \\ & \leq \sum_{\tau=0}^t c^{t-\tau} \sum_{i=1}^N \alpha_i (\mathbb{E}[f_i(y_i(\tau))] - f_i(x^*)) \\ & \leq c^{t+1} \mathbb{E}[V^0] + \sum_{\tau=0}^t c^{t-\tau} \frac{a_M^2 \bar{G}}{2\sigma} \\ & = c^{t+1} \mathbb{E}[V^0] + \frac{a_M^2 \bar{G}}{2\sigma} \frac{1-c^{t+1}}{1-c} \end{aligned} \quad (28)$$

where in the last line we used Lemma 1, thanks to the fact that since by assumption $\alpha_i \leq \frac{Q}{m}$, we have $c \in (0, 1)$. Then, from the convexity of f we have that, at any iteration t ,

$$\begin{aligned} & \sum_{\tau=0}^t c^{t-\tau} a_m (\mathbb{E}[f(\bar{x}(\tau))] - f(x^*)) \\ & \geq \left(a_m \sum_{\tau=0}^t c^{t-\tau} \right) \left(\min_{\tau \leq t} \mathbb{E}[f(\bar{x}(\tau))] - f(x^*) \right) \\ & = \left(a_m \frac{1-c^{t+1}}{1-c} \right) (f_{\text{best}}(\bar{x}(t)) - f(x^*)) \end{aligned} \quad (29)$$

where we used Lemma 1 and the definition of f_{best} . Now, by making some manipulation on the term $\mathbb{E}[f(\bar{x}(\tau))] - f(x^*) = \mathbb{E}[f(\bar{x}(\tau)) - f(x^*)]$, as in [1, Theorem 1] we get

$$\begin{aligned} \mathbb{E}[f(\bar{x}(\tau)) - f(x^*)] & \leq \sum_{i=1}^N \mathbb{E}[(f_i(y_i(\tau)) - f_i(x^*))] \\ & + \sum_{i=1}^N G_i (\mathbb{E}[\|y_i(\tau) - x_i(\tau)\|] + \mathbb{E}[\|x_i(\tau) - \bar{x}(\tau)\|]). \end{aligned} \quad (30)$$

In the case $c \neq \mu_M$, by substituting (30) in (29) and by using (28) and Lemma 3 one has

$$\begin{aligned} & \left(a_m \frac{1-c^{t+1}}{1-c} \right) (f_{\text{best}}(\bar{x}(t)) - f(x^*)) \\ & \leq c^{t+1} \mathbb{E}[V^0] + \frac{a_M^2 \bar{G}}{2\sigma} \frac{1-c^{t+1}}{1-c} \\ & \quad + 3a_m G \left(c^t \left(C + \frac{\bar{R}}{c - \mu_M} \right) + \frac{1-c^t}{1-c} \bar{S} \right). \end{aligned}$$

Now, by dividing both sides by a_m and rearranging the term one has

$$\begin{aligned} & \left(\frac{1-c^{t+1}}{1-c} \right) (f_{\text{best}}(\bar{x}(t)) - f(x^*)) \\ & \leq c^{t+1} \frac{\mathbb{E}[V^0]}{a_m} + c^t 3G \left(C + \frac{\bar{R}}{c - \mu_M} \right) \\ & \quad + \frac{1-c^{t+1}}{1-c} \frac{a_M^2 \bar{G}}{2\sigma a_m} + \frac{1-c^t}{1-c} 3G \bar{S} \\ & \leq c^t \left(\frac{\mathbb{E}[V^0]}{a_m} + 3GC + \frac{3G\bar{R}}{c - \mu_M} \right) \\ & \quad + \frac{1-c^{t+1}}{1-c} \left(\frac{a_M^2 \bar{G}}{2\sigma a_m} + 3G\bar{S} \right) \end{aligned}$$

where in the second line we used the fact that $c \leq 1$. Finally, (26) is obtained by dividing both sides by $\frac{1-c^{t+1}}{1-c}$. The case $c = \mu_m$ can be proven in a similar way. In fact, by using the same arguments as before, we have

$$\begin{aligned} & \left(\frac{1-c^{t+1}}{1-c} \right) (f_{\text{best}}(\bar{x}(t)) - f(x^*)) \\ & \leq c^t \frac{\mathbb{E}[V^0]}{a_m} + c^t 3G \left(C + \frac{t\bar{R}}{c} \right) \\ & \quad + \frac{1-c^{t+1}}{1-c} \left(\frac{a_M^2 \bar{G}}{2\sigma a_m} + 3G\bar{S} \right) \\ & \leq c^t \left(\frac{\mathbb{E}[V^0]}{a_m} + 3GC \right) + tc^t \frac{3G\bar{R}}{c} \\ & \quad + \frac{1-c^{t+1}}{1-c} \left(\frac{a_M^2 \bar{G}}{2\sigma a_m} + 3G\bar{S} \right) \end{aligned}$$

thus leading to (27) by dividing both sides by $\frac{1-c^{t+1}}{1-c}$. \square

Notice that Theorem 1 implies that convergence with a constant error is attained, i.e., define $\tilde{f}^* = f(x^*) + S$, then

$$\lim_{t \rightarrow \infty} f_{\text{best}}(\bar{x}^t) - \tilde{f}^* = 0. \quad (31)$$

Moreover, the convergence rate is linear. In fact, recall that $c \in (0, 1)$. Then, if $c \neq \mu_M$ one has

$$\lim_{t \rightarrow \infty} \frac{f_{\text{best}}(\bar{x}^{t+1}) - \tilde{f}^*}{f_{\text{best}}(\bar{x}^t) - \tilde{f}^*} \leq \lim_{t \rightarrow \infty} \frac{\frac{c^{t+1}}{1-c^{t+2}}}{\frac{c^t}{1-c^{t+1}}} = c,$$

while, if $c = \mu_M$,

$$\lim_{t \rightarrow \infty} \frac{f_{\text{best}}(\bar{x}^{t+1}) - \tilde{f}^*}{f_{\text{best}}(\bar{x}^t) - \tilde{f}^*} \leq \lim_{t \rightarrow \infty} \frac{\frac{c^{t+1}}{1-c^{t+2}} (\bar{\beta} + (t+1)\eta)}{\frac{c^t}{1-c^{t+1}} (\bar{\beta} + t\eta)} = c.$$

Remark 1. Our block-wise algorithm has two main benefits in terms of communication and computation respectively. First, when a limited bandwidth is available in the communication channels, data that exceed the communication bandwidth are transmitted sequentially. For example, if only one block fits the communication channel, our algorithm performs an update at each communication round, while classical ones need B communication rounds per update. Second, in general, solving the minimization problem in (7) on the entire optimization variable or on a single block results in completely different computational times.

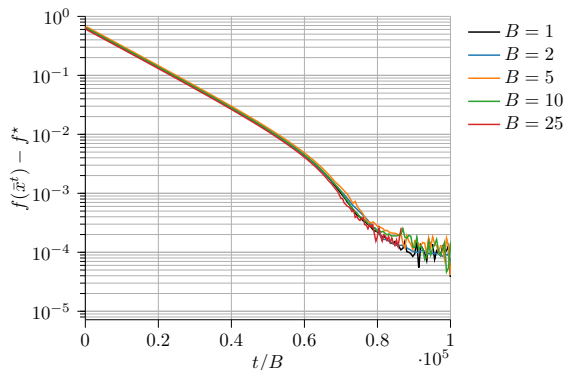


Fig. 1: Numerical example: Evolution of the cost error normalized on the number of blocks.

V. NUMERICAL EXAMPLE

We consider as a numerical example a learning problem in which agents have to classify samples belonging to two clusters. Formally, each agent $i \in \{1, \dots, N\}$ has m_i training samples $q_i^1, \dots, q_i^{m_i} \in \mathbb{R}^d$ each of which has an associated binary label $b_i^r \in \{-1, 1\}$ for all $r \in \{1, \dots, m_i\}$. The goal of the agents is to compute in a distributed way a linear classifier from the training samples, i.e., to find a hyperplane of the form $\{z \in \mathbb{R}^d \mid \langle \theta, z \rangle + \theta_0 = 0\}$, with $\theta \in \mathbb{R}^d$ and $\theta_0 \in \mathbb{R}$, which better separates the training data. For notational convenience, let $x = [\theta^\top, \theta_0]^\top \in \mathbb{R}^{d+1}$ and $\hat{q}_i^r = [(q_i^r)^\top, 1]^\top$. Then, the presented problem can be addressed by solving the following convex optimization problem, in which a regularized Hinge loss is used as cost function,

$$\underset{x \in \mathbb{R}^{d+1}}{\text{minimize}} \quad \sum_{i=1}^N \frac{1}{m_i} \sum_{r=1}^{m_i} \max(0, 1 - b_i^r \langle x, \hat{q}_i^r \rangle) + \frac{\lambda}{2} \|x\|^2,$$

where $\lambda > 0$ is the regularization weight. This problem can be written in the form of (1) by defining $\xi_i^r = (\hat{q}_i^r, b_i^r)$ and

$$\mathbb{E}[h_i(x; \xi_i)] = \frac{1}{m_i} \sum_{r=1}^{m_i} \left(\max(1 - b_i^r \langle x, \hat{q}_i^r \rangle) + \frac{\lambda}{2N} \|x\|^2 \right)$$

for all $i \in \{1, \dots, N\}$. In fact, as long as each data ξ_i^r is uniformly drawn from the dataset, Assumption 1(C) is satisfied. We implemented the algorithm in DISROPT [19] and we tested it in this scenario with $N = 48$ agents, $x \in \mathbb{R}^{50}$ and different number of blocks, namely $B \in \{1, 2, 5, 10, 25\}$. We generated a synthetic dataset composed of 480 points and assigned 10 of them to each agent, i.e., $m_1 = \dots = m_N = 10$. Agents communicate according to a connected graph generated according to an Erdős-Rényi random model with connectivity parameter $p = 0.5$. The corresponding weight matrix is built by using the Metropolis-Hastings rule. Finally, we set $\lambda = 1$, $p_{i,\ell} = 1/B$ for all i and all ℓ , $p_{i,on} = 0.95$ for all i and local (constant) stepsizes α_i randomly chosen according to a normal distribution with mean 0.005 and standard deviation 10^{-4} . The evolution of the cost error adjusted with respect to the number of blocks is reported in Figure 1 for the considered block numbers. The linear convergence rate can be easily appreciated from the figure and confirms the theoretical analysis.

VI. CONCLUSIONS

In this paper, we studied the behavior of the Distributed Block Proximal Method when applied to problems involving (non-smooth) strongly convex functions and when agents in the network employ constant stepsizes. A linear convergence rate (with a constant error) has been obtained in terms of the expected distance from the optimal cost. A numerical example involving a learning problem confirmed the theoretical analysis.

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